

Every Ternary Permutation Constraint Satisfaction Problem Parameterized Above Average Has a Kernel with a Quadratic Number of Variables*

Gregory Gutin¹, Leo van Iersel², Matthias Mnich³, and Anders Yeø¹

¹ Royal Holloway, University of London
Egham, Surrey TW20 0EX, UK
{gutin|anders}@cs.rhul.ac.uk

² University of Canterbury Christchurch, New Zealand
l.j.j.v.iersel@gmail.com

³ Technische Universiteit Eindhoven Eindhoven, The Netherlands
m.mnich@tue.nl

December 20, 2010

Abstract

A ternary Permutation-CSP is specified by a subset Π of the symmetric group \mathcal{S}_3 . An instance of such a problem consists of a set of variables V and a multiset of constraints, which are ordered triples of distinct variables of V . The objective is to find a linear ordering α of V that maximizes the number of triples whose rearrangement (under α) follows a permutation in Π . We prove that every ternary Permutation-CSP parameterized above average has a kernel with a quadratic number of variables.

1 Introduction

For maximization problems whose lower bound on the solution value is a monotonically increasing unbounded function of the instance size, the standard parameterization by solution value is trivially fixed-parameter tractable. (Basic notions on parameterized algorithmics used in this paper are given in Section 2.) Mahajan and Raman [27] were the first to recognize both practical and theoretical importance of parameterizing maximization problems differently: above tight lower bounds. They considered MAX SAT with the tight lower bound $m/2$, where m is the number of clauses, and the problem is to decide whether we can satisfy at least $m/2 + k$ clauses, where k is the parameter. Mahajan and Raman proved that this parameterization of

*Extended abstract of this paper has appeared in the proceedings of the 18th Annual European Symposium on Algorithms, 2010. Part of this research has been supported by the EPSRC, grant EP/E034985/1, the Netherlands Organisation for Scientific Research (NWO), grant 639.033.403, and the Allan Wilson Centre for Molecular Ecology and Evolution.

MAX SAT is fixed-parameter tractable by obtaining a kernel with $O(k)$ variables. Despite clear importance of parameterizations above tight lower bounds, until recently only a few sporadic non-trivial results on the topic were obtained [19, 22, 23, 27, 31].

Massive interest in parameterizations above tight lower bounds came with the paper of Mahajan et al. [28], who stated several questions on fixed-parameter tractability of maximization problems parameterized above tight lower bounds, some of which are still open. Several of those questions were answered by newly-developed methods [1, 9, 10, 20, 21], using algebraic, probabilistic and harmonic analysis tools. In particular, a probabilistic approach allowed Gutin et al. [20] to prove the existence of a quadratic kernel for the parameterized BETWEENNESS ABOVE AVERAGE (BETWEENNESS-AA) problem, thus, answering an open question of Benny Chor [29].

BETWEENNESS is just one representative of a rich family of *ternary Permutation Constraint Satisfaction Problems (CSPs)*. A ternary Permutation-CSP is specified by a subset Π of the symmetric group \mathcal{S}_3 . An instance of such a problem consists of a set of variables V and a multiset of constraints, which are ordered triples of distinct variables of V . The objective is to find a linear ordering α of V that maximizes the number of triples whose rearrangement (under α) follows a permutation in Π . Important special cases are BETWEENNESS [6, 16, 20, 30] and CIRCULAR ORDERING [15, 17], which find applications in circuit design and computational biology [8, 30], and in qualitative spatial reasoning [26], respectively.

In this paper, we prove that every ternary Permutation-CSP has a kernel with a quadratic number of variables, when parameterized above average (AA), which is a tight lower bound. This result is obtained by first reducing all the problems to just one, LINEAR ORDERING-AA, then showing that LINEAR ORDERING-AA has a kernel with a quadratic number of variables and constraints and, thus, concluding that there is a bikernel with a quadratic number of variables from each of the problems AA to LINEAR ORDERING-AA. Using the last result, we prove that there is a bikernel with a quadratic number of variables from every ternary Permutation-CSP to most ternary Permutation-CSPs. This implies the existence of kernels with a quadratic number of variables for most ternary Permutation-CSPs. The remaining ternary Permutation-CSPs are proved to be equivalent to ACYCLIC SUBDIGRAPH-AA (a *binary Permutation-CSP* defined in Section 6) and since ACYCLIC SUBDIGRAPH-AA, as shown in [21], has a kernel with a quadratic number of variables, the remaining ternary Permutation-CSPs have a kernel with a quadratic number of variables.

The most difficult part of this set of arguments is the proof that LINEAR ORDERING-AA has a kernel with a quadratic number of variables and constraints. We can show that if we want to prove this in a similar way as for BETWEENNESS-AA (that is, eliminate all instances of LINEAR ORDERING-AA whose optimal solution coincides with the lower bound) we need an infinite number of reduction rules, see Section 8 for details. So, determining fixed-parameter tractability of LINEAR ORDERING-AA turns out to be much harder than that for BETWEENNESS-AA. Fortunately, we found a nontrivial way of reducing LINEAR ORDERING-AA to a combination of BETWEENNESS-AA and ACYCLIC SUBDIGRAPH-AA. Using further probabilistic and deterministic arguments for the mixed problem, we prove that LINEAR ORDERING-AA has a kernel with a quadratic number of variables and constraints.

The rest of the paper is organized as follows. In Section 2, we give some basic notions on parameterized algorithms and complexity. In Section 3, we define and

discuss ternary Permutation-CSPs; we also reduce all nontrivial ternary Permutation-CSPs AA to LINEAR ORDERING-AA. Some NP-hardness results of this section are proved in Section 4. In Section 5, we describe probabilistic and harmonic analysis tools used in the paper. In Section 6, we obtain some results on BETWEENNESS-AA and ACYCLIC SUBDIGRAPH-AA needed in the following section, where we prove that LINEAR ORDERING-AA has a quadratic kernel. In Section 7, we also prove our main result, Theorem 5, that every ternary Permutation-CSs parameterized above average has a kernel with a quadratic number of variables. In Section 8, we show that LINEAR ORDERING-AA has an infinite number of natural reduction rules. Finally, in Section 9 we state and discuss open problems for further research.

2 Basics on Parameterized Algorithmics

Parameterized complexity theory is a multivariate framework for a refined analysis of hard (NP-hard) problems, which was introduced by Downey and Fellows in a series of ground breaking papers in the 1990s [11, 12]. A *parameterized problem* is a subset $L \subseteq \Sigma^* \times \mathbb{N}$ over a finite alphabet Σ ; L is *fixed-parameter tractable* if the membership of an instance (I, k) in $\Sigma^* \times \mathbb{N}$ can be decided in time $f(k) \cdot |I|^{O(1)}$ where f is a function of the *parameter* k only [13, 14, 29]. (We would like $f(k)$ to grow as slowly as possible.)

Given a pair L, L' of parameterized problems, a *bikernelization from L to L'* is a polynomial-time algorithm that maps an instance (x, k) to an instance (x', k') (the *bikernel*) such that (i) $(x, k) \in L$ if and only if $(x', k') \in L'$, (ii) $k' \leq h(k)$, and (iii) $|x'| \leq g(k)$ for some functions h and g . The function $g(k)$ is called the *size* of the bikernel. A *kernelization* of a parameterized problem L is simply a bikernelization from L to itself and a *bikernel* is a *kernel* when $L = L'$.

The notion of a bikernelization was introduced by Alon et al. [1], who observed that a decidable parameterized problem L is fixed-parameter tractable if and only if it admits a bikernelization to a decidable parameterized problem L' . Not every fixed-parameter tractable problem has a kernel of polynomial size unless $\text{NP} \subseteq \text{coNP/poly}$ [3, 4, 2]; low degree polynomial size kernels are of main interest due to applications.

3 Permutation CSPs Parameterized Above Average

Let V be a set of n variables. A *linear ordering of V* is a bijection $\alpha : V \rightarrow [n]$, where $[n] = \{1, 2, \dots, n\}$. The symmetric group on three elements is $\mathcal{S}_3 = \{(123), (132), (213), (231), (312), (321)\}$. A *constraint set over V* is a multiset \mathcal{C} of *constraints*, which are permutations of three distinct elements of V . For each subset $\Pi \subseteq \mathcal{S}_3$ and a linear ordering α of V , a constraint $(v_1, v_2, v_3) \in \mathcal{C}$ is *Π -satisfied by α* if there is a permutation $\pi \in \Pi$ such that $\alpha(v_{\pi(1)}) < \alpha(v_{\pi(2)}) < \alpha(v_{\pi(3)})$. If Π is fixed, we will simply say that $(v_1, v_2, v_3) \in \mathcal{C}$ is *satisfied by α* .

For each subset $\Pi \subseteq \mathcal{S}_3$, the problem Π -CSP is to decide whether for a given pair (V, \mathcal{C}) of variables and constraints there is a linear ordering α of V that Π -satisfies all constraints in \mathcal{C} . A complete dichotomy of the Π -CSP problems with respect to their computational complexity was given by Guttman and Maucher [25]. For

$\Pi \subseteq \mathcal{S}_3$	Common Problem Name	Complexity to Satisfy All Constraints
$\Pi_0 = \{(123)\}$	LINEAR ORDERING	polynomial
$\Pi_1 = \{(123), (132)\}$		polynomial
$\Pi_2 = \{(123), (213), (231)\}$		polynomial
$\Pi_3 = \{(132), (231), (312), (321)\}$		polynomial
$\Pi_4 = \{(123), (231)\}$		NP-complete
$\Pi_5 = \{(123), (321)\}$	BETWEENNESS	NP-complete
$\Pi_6 = \{(123), (132), (231)\}$		NP-complete
$\Pi_7 = \{(123), (231), (312)\}$	CIRCULAR ORDERING	NP-complete
$\Pi_8 = \mathcal{S}_3 \setminus \{(123), (231)\}$		NP-complete
$\Pi_9 = \mathcal{S}_3 \setminus \{(123), (321)\}$	NON-BETWEENNESS	NP-complete
$\Pi_{10} = \mathcal{S}_3 \setminus \{(123)\}$		NP-complete

Table 1: Ternary Permutation-CSPs (after symmetry considerations)

that, they reduced $2^{|\mathcal{S}_3|} = 64$ problems by two types of symmetry. First, two problems differing just by a consistent renaming of the elements of their permutations are of the same complexity. Second, two problems differing just by reversing their permutations are of the same complexity. The symmetric reductions leave 13 problems Π_i -CSP, $i = 0, 1, \dots, 12$, whose time complexity is polynomial for $\Pi_{11} = \emptyset$ and $\Pi_{12} = \mathcal{S}_3$ and was otherwise established by Guttman and Maucher [25], see Table 1.

The maximization version of Π_i -CSP is the problem MAX- Π_i -CSP of finding a linear ordering α of V that Π_i -satisfies a maximum number of constraints in \mathcal{C} . Clearly, for $i = 4, \dots, 10$ the problem MAX- Π_i -CSP is NP-hard. In Section 4 we prove that MAX- Π_i -CSP is NP-hard also for $i = 0, 1, 2, 3$.

Now observe that given a variable set V and a constraint multiset \mathcal{C} over V , for a random linear ordering α of V , the probability of a constraint in \mathcal{C} being Π -satisfied by α equals $\frac{|\Pi|}{6}$. Hence, the expected number of satisfied constraints from \mathcal{C} is $\frac{|\Pi|}{6}|\mathcal{C}|$, and thus there is a linear ordering α of V satisfying at least $\frac{|\Pi|}{6}|\mathcal{C}|$ constraints (and this bound is tight). A derandomization argument leads to $\frac{|\Pi|}{6}$ -approximation algorithms for the problems MAX- Π_i -CSP [6]. No better constant factor approximation is possible assuming the Unique Games Conjecture [6].

We study the parameterization of MAX- Π_i -CSP above tight lower bound:

Π -ABOVE AVERAGE (Π -AA)

Input: A finite set V of variables, a multiset \mathcal{C} of ordered triples of distinct variables from V and an integer $k \geq 0$.

Parameter: k .

Question: Is there a linear ordering α of V such that at least $\frac{|\Pi|}{6}|\mathcal{C}| + k$ constraints of \mathcal{C} are Π -satisfied by α ?

For example, choose $\Pi = \{(123), (321)\}$ for BETWEENNESS-AA. We will call Π_0 -

AA the LINEAR ORDERING-AA problem.

Let Π be a subset of \mathcal{S}_3 . Clearly, if Π is the empty set or equal to \mathcal{S}_3 then the corresponding problem Π -AA can be solved in polynomial time. The following simple result allows us to study the Π -AA problems using Π_0 -AA.

Proposition 1. *Let Π be a subset of \mathcal{S}_3 such that $\Pi \notin \{\emptyset, \mathcal{S}_3\}$. There is a polynomial time transformation f from Π -AA to Π_0 -AA such that an instance (V, \mathcal{C}, k) of Π -AA is a “yes”-instance if and only if $(V, \mathcal{C}_0, k) = f(V, \mathcal{C}, k)$ is a “yes”-instance of Π_0 -AA.*

Proof. From an instance (V, \mathcal{C}, k) of Π -AA, construct an instance (V, \mathcal{C}_0, k) of Π_0 -AA as follows. For each triple $(v_1, v_2, v_3) \in \mathcal{C}$, add $|\Pi|$ triples $(v_{\pi(1)}, v_{\pi(2)}, v_{\pi(3)})$, $\pi \in \Pi$, to \mathcal{C}_0 .

Observe that a triple $(v_1, v_2, v_3) \in \mathcal{C}$ is Π -satisfied if and only if exactly one of the triples $(v_{\pi(1)}, v_{\pi(2)}, v_{\pi(3)})$, $\pi \in \Pi$, is Π_0 -satisfied. Thus, $\frac{|\Pi|}{6}|\mathcal{C}| + k$ constraints from \mathcal{C} are Π -satisfied if and only if the same number of constraints from \mathcal{C}_0 are Π_0 -satisfied. It remains to observe that $\frac{|\Pi|}{6}|\mathcal{C}| + k = \frac{1}{6}|\mathcal{C}_0| + k$ as $|\mathcal{C}_0| = |\Pi| \cdot |\mathcal{C}|$. \square

For a variable set V , a constraint multiset \mathcal{C} over V and a linear ordering α of V , the α -deviation of (V, \mathcal{C}) is the number $\text{dev}(V, \mathcal{C}, \alpha)$ of constraints of \mathcal{C} that are Π -satisfied by α minus $\frac{|\Pi|}{6}|\mathcal{C}|$. The *maximum deviation* of (V, \mathcal{C}) , denoted $\text{dev}(V, \mathcal{C})$, is the maximum of $\text{dev}(V, \mathcal{C}, \alpha)$ over all linear orderings α of V . Now the problem Π -AA can be reformulated as the problem of deciding whether $\text{dev}(V, \mathcal{C}) \geq k$.

4 NP-hardness of MAX- Π_i -CSP for $i = 0, 1, 2, 3$

The problem ACYCLIC SUBDIGRAPH is, given a directed multigraph D and an integer $k > 0$, to decide whether D contains an acyclic subdigraph with at least k arcs. ACYCLIC SUBDIGRAPH can be reformulated as a problem of verifying whether V has a linear ordering α in which at least k arcs are *satisfied*, i.e., for each such arc (u, v) we have $\alpha(u) < \alpha(v)$. It is well-known that ACYCLIC SUBDIGRAPH is NP-complete.

Theorem 1. *For $i = 0, 1, 2, 3$, problem MAX- Π_i -CSP from Table 1 is NP-hard.*

Proof. We will consider the four cases one by one.

$i = 0$: Proposition 1 implies, in particular, that MAX-BETWEENNESS can be reduced to MAX- Π_0 -CSP. Thus, MAX- Π_0 -CSP is NP-hard.

$i = 1$: Denote constraints of MAX- Π_1 -CSP by $(u < \min\{v, w\})$. Such a constraint is Π_1 -satisfied by a linear ordering α of $\{u, v, w\}$ if and only if $\alpha(u) < \min\{\alpha(v), \alpha(w)\}$. From an instance $(D = (U, A), k)$ of ACYCLIC SUBDIGRAPH, we construct an instance (V, \mathcal{C}, k) of (a decision version of) MAX- Π_1 -CSP by setting $V = U \cup \{z\}$ and, for each arc $(u, v) \in A$, adding $(u < \min\{v, z\})$ to \mathcal{C} . Observe that, without loss of generality, an optimal linear ordering of (V, \mathcal{C}, k) has z at the end as if it does not then moving z to the end does not falsify any constraints. Therefore (u, v) is satisfied in D if and only if $(u < \min\{v, z\})$ is Π_1 -satisfied in (V, \mathcal{C}, k) .

$i = 2$: Denote constraints of MAX- Π_2 -CSP by $(u, v < w)$. Such a constraint is Π_2 -satisfied by a linear ordering α of $\{u, v, w\}$ if and only if $\alpha(v) < \alpha(w)$. From

an instance $(D = (U, A), k)$ of ACYCLIC SUBDIGRAPH, we construct an instance (V, \mathcal{C}, k) of (a decision version of) MAX- Π_2 -CSP by setting $V = U \cup \{z\}$ and, for each arc $(v, w) \in A$, adding constraint $(z, v < w)$ to the constraint set \mathcal{C} . Observe that D has a set of k arcs that form an acyclic subdigraph if and only if there are k constraints in \mathcal{C} that can be Π_2 -satisfied by a linear ordering of V . Thus, we have reduced ACYCLIC SUBDIGRAPH to MAX- Π_2 -CSP, implying that MAX- Π_2 -CSP is NP-hard.

$i = 3$: Let us denote a constraint in MAX- Π_3 -CSP by $(\max\{u, v\} \not< w)$. This constraint is Π_3 -satisfied by a linear ordering α if and only if w is not the last element among u, v, w in α . Now consider an instance (V, \mathcal{C}_1, k) of MAX- Π_1 -CSP, which we have shown to be NP-hard. For each constraint $(u < \min\{v, w\})$ in \mathcal{C}_1 add $(\max\{u, v\} \not< w)$ and $(\max\{u, w\} \not< v)$ to \mathcal{C}_3 . Now we will show that (V, \mathcal{C}_1, k) is a “yes”-instance if and only if $(V, \mathcal{C}_3, |\mathcal{C}_1| + k)$ is a “yes”-instance of MAX- Π_3 -CSP. Let α be any linear ordering of V and let α' be the reverse ordering. Note that $(u < \min\{v, w\})$ is Π_1 -satisfied by α if and only if both $(\max\{u, v\} \not< w)$ and $(\max\{u, w\} \not< v)$ are Π_3 -satisfied by α' . Furthermore one of $(\max\{u, v\} \not< w)$ and $(\max\{u, w\} \not< v)$ is always Π_3 -satisfied. Therefore, at least k constraints of \mathcal{C}_1 are Π_1 -satisfied by α if and only if at least $2k + (|\mathcal{C}_1| - k)$ constraints of \mathcal{C}_3 are Π_3 -satisfied by α' . So, we have reduced MAX- Π_1 -CSP to MAX- Π_3 -CSP, and we are done. □

5 Probabilistic and Harmonic Analysis Tools

We build on the probabilistic *Strictly Above Expectation* method by Gutin et al. [21] to prove non-trivial lower bounds on the minimum fraction of satisfiable constraints in instances belonging to a restricted subclass. For such an instance with parameter k , we introduce a random variable X such that the instance is a “yes”-instance if and only if X takes with positive probability a value greater than or equal to k . If X happens to be a symmetric random variable with finite second moment then $\mathbb{P}(X \geq \sqrt{\mathbb{E}[X^2]}) > 0$; it hence suffices to prove $\mathbb{E}[X^2] = h(k)$ for some monotonically increasing unbounded function h . (Here, $\mathbb{P}(\cdot)$ and $\mathbb{E}[\cdot]$ denote probability and expectation, respectively.) If X is not symmetric then the following lemma can be used instead.

Lemma 1 (Alon et al. [1]). *Let X be a real random variable and suppose that its first, second and fourth moments satisfy $\mathbb{E}[X] = 0$, $\mathbb{E}[X^2] = \sigma^2 > 0$ and $\mathbb{E}[X^4] \leq c\sigma^4$, respectively, for some constant c . Then $\mathbb{P}(X > \frac{\sigma}{2\sqrt{c}}) > 0$.*

We combine this result with the following result from harmonic analysis.

Lemma 2 (Hypercontractive Inequality [5, 18]). *Let $f = f(x_1, \dots, x_n)$ be a polynomial of degree r in n variables x_1, \dots, x_n with domain $\{-1, 1\}$. Define a random variable X by choosing a vector $(\epsilon_1, \dots, \epsilon_n) \in \{-1, 1\}^n$ uniformly at random and setting $X = f(\epsilon_1, \dots, \epsilon_n)$. Then $\mathbb{E}[X^4] \leq 9^r \mathbb{E}[X^2]^2$.*

$ \{\phi(u), \phi(v), \phi(w)\} $	Relation	Value of Y_p	Prob.
1	$\phi(u) = \phi(v) = \phi(w)$	0	1/16
2	$\phi(v) \neq \phi(u) = \phi(w)$	-1/3	3/16
2	$\phi(v) \in \{\phi(u), \phi(w)\}$	1/6	6/16
3	$\phi(v)$ is between $\phi(u)$ and $\phi(w)$	2/3	2/16
3	$\phi(v)$ is not between $\phi(u)$ and $\phi(w)$	-1/3	4/16

Table 2: Distribution of Y_p for constraint $B_p = (v, \{u, w\})$.

6 Betweenness and Acyclic Subdigraph Problems

Let u, v, w be variables. We denote a betweenness constraint “ v is between u and w ” by $(v, \{u, w\})$, and call a 3-set S of betweenness constraints over $\{u, v, w\}$ *complete* if $S = \{(u, \{v, w\}), (v, \{u, w\}), (w, \{u, v\})\}$. Since every linear ordering of $\{u, v, w\}$ satisfies exactly one constraint in S , we obtain the following reduction.

Lemma 3. *Let (V, \mathcal{B}) be an instance of BETWEENNESS and let α be a linear ordering of V . Let \mathcal{B}' be the set of constraints obtained from \mathcal{B} by deleting all complete subsets. Then $\text{dev}(V, \mathcal{B}, \alpha) = \text{dev}(V, \mathcal{B}', \alpha)$.*

An instance of BETWEENNESS without complete subsets of constraints is called *reduced*.

Let (V, \mathcal{B}) be an instance of BETWEENNESS, with $\mathcal{B} = \{B_1, \dots, B_m\}$, and let ϕ be a fixed function from V to $\{0, 1, 2, 3\}$. A linear ordering α of V is called ϕ -*compatible* if for each pair $u, v \in V$ with $\alpha(u) < \alpha(v)$ it holds $\phi(u) \leq \phi(v)$. For a random ϕ -compatible linear ordering π of V , define a binary random variable y_p that takes value one if and only if $B_p \in \mathcal{B}$ is satisfied by π (if B_p is falsified by π , then $y_p = 0$). Let $Y_p = \mathbb{E}[y_p] - 1/3$ for each $p \in [m]$, and let $Y = \sum_{p=1}^m Y_p$.

Now let ϕ be a random function from V to $\{0, 1, 2, 3\}$. Then Y, Y_1, \dots, Y_m are random variables. For a constraint $B_p = (v, \{u, w\})$, the distribution of Y_p as it is given in Table 2 implies that $\mathbb{E}[Y_p] = 0$. Thus, by linearity of expectation, $\mathbb{E}[Y] = 0$.

The following lemma was proved by Gutin et al. [20] for BETWEENNESS in which \mathcal{B} is a set, not a multiset, but a simple modification of its proof gives us the following:

Lemma 4. *For a reduced instance (V, \mathcal{B}) of BETWEENNESS, $\mathbb{E}[Y^2] \geq \frac{11}{768}m$.*

Proof. Observe that $\mathbb{E}[Y^2] = \sum_{l=1}^m \mathbb{E}[Y_l^2] + \sum_{1 \leq l \neq l' \leq m} \mathbb{E}[Y_l Y_{l'}]$. Using Table 2, it is easy to see that $\sum_{l=1}^m \mathbb{E}[Y_l^2] = \frac{88}{768}m$.

Let $U = \{(l, l') : B_l, B_{l'} \in \mathcal{B}, l \neq l'\}$ be the set of all ordered index pairs corresponding to distinct constraints in \mathcal{B} . Let $U^* = \{(l, l') \in U : \text{vars}(B_l) = \text{vars}(B_{l'}), B_l \neq B_{l'}\}$ and $U^{**} = \{(l, l') \in U : B_l = B_{l'}\}$, where $\text{vars}(B_l)$ and $\text{vars}(B_{l'})$ are the sets of variables of B_l and $B_{l'}$, respectively. Taking into consideration that $|U^*| \leq m$ and $|U^{**}| \geq 0$, similarly to [20], we obtain that

$$\sum_{(l, l') \in U} \mathbb{E}[Y_l Y_{l'}] \geq -\frac{66}{768}m - \frac{11}{768}|U^*| + \frac{22}{768}|U^{**}| \geq -\frac{66}{768}m - \frac{11}{768}m = -\frac{77}{768}m.$$

Relation between $\phi(u)$ and $\phi(v)$	Value of X_p	Prob.
$\phi(u) = \phi(v)$	0	1/4
$\phi(u) < \phi(v)$	1/2	3/8
$\phi(u) > \phi(v)$	-1/2	3/8

Table 3: Distribution of X_p for an arc (u, v) .

Combining this with $\sum_{l=1}^m \mathbb{E}[Y_l^2] = \frac{88}{768}m$, we get $\mathbb{E}[Y^2] \geq \frac{11}{768}m$. \square

Recall that in the ACYCLIC SUBDIGRAPH problem we are given a directed multigraph $D = (U, A)$, with parallel arcs allowed, and ask for a linear ordering π of V which maximizes the number of satisfied arcs, where an arc $(u, v) \in A$ is *satisfied* by π if $\pi(u) < \pi(v)$. If π is a uniformly-at-random linear ordering of V then the probability of an arc of D being satisfied is $1/2$. Thus, there is a linear ordering π of V in which the number of satisfied arcs is at least $|A|/2$. We therefore define, for a digraph $D = (U, A)$ and a linear ordering π of U , the π -*deviation* of D as the number of arcs satisfied by π minus $|A|/2$, and denote it by $\text{dev}(V, A, \pi)$. In the ACYCLIC SUBDIGRAPH-AA problem we are given a directed multigraph $D = (U, A)$ and asked to decide whether there is a linear ordering π of U with π -*deviation* at least k , where k is a parameter.

As every linear ordering of U satisfies exactly one of two mutually opposite arcs (u, v) and (v, u) , we obtain the following reduction.

Lemma 5. *Let $D = (U, A)$ be a directed multigraph and let π be a linear ordering of V . Let A' be the set of arcs obtained from A by deleting all pairs of mutually opposite arcs. Then $\text{dev}(V, A, \pi) = \text{dev}(V, A', \pi)$.*

A directed multigraph without mutually opposite arcs is called *reduced*.

Let $D = (U, A)$ be a directed multigraph with $A = \{a_1, \dots, a_m\}$ as multiset of arcs, and let ϕ be a fixed function from U to $\{0, 1, 2, 3\}$. For a random ϕ -compatible linear ordering π of U , define a binary random variable x_p that takes value one if and only if a_p is satisfied by π . Let $X_p = \mathbb{E}[x_p] - 1/2$ for each $p \in [m]$ and let $X = \sum_{p=1}^m X_p$.

Now let ϕ be a random function from U to $\{0, 1, 2, 3\}$. Then X, X_1, \dots, X_m are random variables. For an arc (u, v) , the distribution of X_p as it is given in Table 3 implies that $\mathbb{E}[X_p] = 0$. Thus, by linearity of expectation, $\mathbb{E}[X] = 0$.

We have the following analogue of Lemma 4.

Lemma 6. *For reduced directed multigraphs D it holds that $\mathbb{E}[X^2] \geq \frac{1}{32}m$.*

Proof. We write $\mathbb{E}[X^2]$ as the sum

$$\mathbb{E}[X^2] = \sum_{p=1}^m \mathbb{E}[X_p^2] + \sum_{1 \leq p \neq p' \leq m} \mathbb{E}[X_p X_{p'}]. \quad (1)$$

From Table 3 it follows that $\mathbb{E}[X_p^2] = \frac{3}{16}$, and hence it remains to bound the second sum in (1). Consider any ordered pair $(a_p, a_{p'})$ of distinct arcs in D . If a_p and $a_{p'}$ are

vertex-disjoint, then clearly $\mathbb{E}[X_p X_{p'}] = 0$. If a_p and $a_{p'}$ have vertices in common, we define

$$\begin{aligned} S_1(u) &= \{(p, p') \mid a_p = (u, x), a_{p'} = (u, y), x, y \in V\} \\ &\quad \cup \{(p, p') \mid a_p = (x, u), a_{p'} = (y, u), x, y \in V\} \\ S_2(u) &= \{(p, p') \mid a_p = (u, x), a_{p'} = (y, u), x, y \in V\} \\ &\quad \cup \{(p, p') \mid a_p = (x, u), a_{p'} = (u, y), x, y \in V\} \\ S_3(u, v) &= \{(p, p') \mid a_p = (u, v), a_{p'} = (u, v)\} . \end{aligned}$$

By setting $l(u) = |\{a \in A : a = (u, y), y \in V\}|$ and $r(u) = |\{a \in A : a = (x, u), x \in V\}|$ it follows that

$$\begin{aligned} |S_1(u)| &= l(u)(l(u) - 1) + r(u)(r(u) - 1), \\ |S_2(u)| &= 2l(u)r(u). \end{aligned}$$

Consider a pair $(p, p') \in S_1(u)$, with say $a_p = (u, x), a_{p'} = (u, y)$. It is easy to calculate that out of the 64 functions $\phi : \{u, x, y\} \rightarrow \{0, 1, 2, 3\}$, there are 14 functions in which $\phi(u) < \phi(x)$ and $\phi(u) < \phi(y)$. Symmetrically, there are 14 functions ϕ in which $\phi(u) > \phi(x)$ and $\phi(u) > \phi(y)$. In both cases, $X_p X_{p'} = \frac{1}{4}$, by Table 3. Similarly, there are 4 functions ϕ in which $\phi(u) < \phi(x)$ and $\phi(u) > \phi(y)$, and 4 functions ϕ in which $\phi(u) > \phi(x)$ and $\phi(u) < \phi(y)$; in both cases $X_p X_{p'} = -\frac{1}{4}$. For all other functions ϕ we have that $X_p X_{p'} = 0$, and thus it follows that $\mathbb{E}[X_p X_{p'}] = \frac{5}{64}$ for each pair of arcs $(a_p, a_{p'})$ in $S_1(u)$.

Similarly, for each pair $(p, p') \in S_2(u)$ it holds that $\mathbb{E}[X_p X_{p'}] = -\frac{5}{64}$, and for each pair $(p, p') \in S_3(u, v)$ it holds that $\mathbb{E}[X_p X_{p'}] = \mathbb{E}[X_p^2] = \frac{3}{16}$.

Hence,

$$\sum_{1 \leq p \neq p' \leq m} \mathbb{E}[X_p X_{p'}] = \sum_{u \in V} \frac{5}{64} |S_1(u)| - \frac{5}{64} |S_2(u)| + \sum_{u, v \in V} w' |S_3(u, v)|,$$

with $\frac{5}{64} + \frac{5}{64} + w' = \frac{3}{16}$, because $S_3(u, v) = S_1(u) \cap S_1(v)$. Thus, $w' = \frac{1}{32}$, and we obtain

$$\begin{aligned} &\sum_{1 \leq p \neq p' \leq m} \mathbb{E}[X_p X_{p'}] \\ &= \frac{5}{64} \sum_{u \in V} l(u)(l(u) - 1) + r(u)(r(u) - 1) - 2l(u)r(u) + \sum_{u, v \in V} \frac{1}{32} |S_3(u, v)| \\ &= \frac{5}{64} \sum_{u \in V} (l(u) - r(u))^2 - l(u) - r(u) + \sum_{u, v \in V} \frac{1}{32} |S_3(u, v)| \\ &\geq -\frac{5}{64} \sum_{u \in V} l(u) + r(u) = -\frac{10}{64} m, \end{aligned}$$

because each arc contributes exactly one to $\sum_{u \in V} l(u)$ and one to $\sum_{u \in V} r(u)$. We conclude that $\mathbb{E}[X^2] \geq \frac{3}{16} m - \frac{10}{64} m = \frac{1}{32} m$. \square

The following theorem was proved in [21].

Theorem 2. ACYCLIC SUBDIGRAPH-AA has a kernel with a quadratic number of vertices and arcs.

linear ordering π of $\{u, v, w\}$	constraints satisfied by π
uvw	$(u, v), (v, w), (v, \{u, w\})$
uww	(u, v)
wuv	(u, v)
vuw	(v, w)
vwu	(v, w)
wvu	$(v, \{u, w\})$

Table 4: Constraints satisfied by π .

7 Kernels for Π -AA Problems

We start from the following key construction of this paper. With an instance (V, \mathcal{C}) of LINEAR ORDERING, we associate an instance (V, \mathcal{B}) of BETWEENNESS and two instances (V, A') and (V, A'') of ACYCLIC SUBDIGRAPH as follows: If $C_p = (u, v, w) \in \mathcal{C}$, then $B_p = (v, \{u, w\}) \in \mathcal{B}$, $a'_p = (u, v) \in A'$, and $a''_p = (v, w) \in A''$.

Lemma 7. *Let (V, \mathcal{C}, k) be an instance of LINEAR ORDERING-AA and let α be a linear ordering of V . Then*

$$\text{dev}(V, \mathcal{C}, \alpha) = \frac{1}{2} [\text{dev}(V, A', \alpha) + \text{dev}(V, A'', \alpha) + \text{dev}(V, \mathcal{B}, \alpha)].$$

Proof. For each constraint $C_p = (u, v, w) \in \mathcal{C}$, define a binary variable \hat{x}'_p that takes value one if and only if a'_p is satisfied by α . Similarly, define binary variables \hat{x}''_p for arc a''_p , \hat{y}_p for constraint B_p and \hat{z}_p for constraint C_p . To show the lemma it suffices to prove that for each constraint $C_p \in \mathcal{C}$ and every linear ordering π of $\{x, y, z\}$ it holds that

$$\text{dev}(V, \{C_p\}, \pi) = \frac{1}{2} [\text{dev}(V, \{a'_p\}, \pi) + \text{dev}(V, \{a''_p\}, \pi) + \text{dev}(V, \{B_p\}, \pi)],$$

where $\text{dev}(V, \{C_p\}, \pi) = \hat{z}_p - 1/6$, $\text{dev}(V, \{a'_p\}, \pi) = \hat{x}'_p - 1/2$, $\text{dev}(V, \{a''_p\}, \pi) = \hat{x}''_p - 1/2$ and $\text{dev}(V, \{B_p\}, \pi) = \hat{y}_p - 1/3$. Thus, it suffices to prove that $\hat{z}_p = (\hat{x}'_p + \hat{x}''_p + \hat{y}_p - 1)/2$. But this expression holds, as can be seen from Table 4: if C_p is satisfied by π then all three constraints a'_p, a''_p, B_p are satisfied by π , whereas if C_p is not satisfied by π then exactly one of the three constraints a'_p, a''_p, B_p is satisfied by π . \square

Let (V, \mathcal{C}, k) be an instance of LINEAR ORDERING-AA, and let ϕ be a function from V to $\{0, 1, 2, 3\}$. For a random ϕ -compatible linear ordering π of V , define a binary random variable z_p that takes value one if and only if C_p is satisfied by π . Let $Z_p = \mathbb{E}[z_p] - 1/6$ for each $p \in [m]$, and let $Z = \sum_{p=1}^m Z_p$.

Lemma 8. *If $Z \geq k$ then (V, \mathcal{C}, k) is a “yes”-instance of LINEAR ORDERING-AA.*

Proof. By linearity of expectation, $Z \geq k$ implies $\mathbb{E}[\sum_{p=1}^m z_p] \geq m/6 + k$. Thus, if $Z \geq k$ then there is a ϕ -compatible permutation π that satisfies at least $m/6 + k$ constraints. \square

Relation between $\phi(u)$, $\phi(v)$ and $\phi(w)$	X'_p	X''_p	Y_p	Z_p
$\phi(u) = \phi(v) = \phi(w)$	0	0	0	0
$\phi(v) < \phi(u) = \phi(w)$	-1/2	1/2	-1/3	-1/6
$\phi(v) > \phi(u) = \phi(w)$	1/2	-1/2	-1/3	-1/6
$\phi(v) = \phi(u) < \phi(w)$	0	1/2	1/6	1/3
$\phi(v) = \phi(u) > \phi(w)$	0	-1/2	1/6	-1/6
$\phi(u) < \phi(v) = \phi(w)$	1/2	0	1/6	1/3
$\phi(u) > \phi(v) = \phi(w)$	-1/2	0	1/6	-1/6
$\phi(u) < \phi(v) < \phi(w)$	1/2	1/2	2/3	5/6
$\phi(u) < \phi(w) < \phi(v)$	1/2	-1/2	-1/3	-1/6
$\phi(v) < \phi(u) < \phi(w)$	-1/2	1/2	-1/3	-1/6
$\phi(v) < \phi(w) < \phi(u)$	-1/2	1/2	-1/3	-1/6
$\phi(w) < \phi(u) < \phi(v)$	1/2	-1/2	-1/3	-1/6
$\phi(w) < \phi(v) < \phi(u)$	-1/2	-1/2	2/3	-1/6

Table 5: Values of X'_p, X''_p, Y_p, Z_p .

Fix a function $\phi : V \rightarrow \{0, 1, 2, 3\}$ and assign variables Y_p, X'_p, X''_p , respectively, to the three instances of BETWEENNESS and ACYCLIC SUBDIGRAPH above.

Lemma 9. *For each $p \in [m]$, we have $Z_p = \frac{1}{2} [X'_p + X''_p + Y_p]$.*

Proof. Let $C_p = (u, v, w) \in \mathcal{C}$. Table 5 shows the values of X'_p, X''_p, Y_p, Z_p for some relations between $\phi(u), \phi(v)$ and $\phi(w)$. The values of X'_p, X''_p and Y_p can be computed using Tables 2 and 3. In all cases of Table 5 it holds $Z_p = \frac{1}{2}(X'_p + X''_p + Y_p)$. Thus, $Z_p = \frac{1}{2}[X'_p + X''_p + Y_p]$ for each possible relation between $\phi(u), \phi(v)$ and $\phi(w)$. \square

Let $X = \sum_{p=1}^m [X'_p + X''_p]$, let $Y = \sum_{p=1}^m Y_p$ and let ϕ be a random function from V to $\{0, 1, 2, 3\}$. Then $X, X'_1, \dots, X'_m, X''_1, \dots, X''_m, Y, Y_1, \dots, Y_m, Z, Z_1, \dots, Z_m$ are random variables. From $\mathbb{E}[X'] = \mathbb{E}[X''] = \mathbb{E}[Y] = 0$ it follows that $\mathbb{E}[Z] = 0$.

We will be able to use Lemma 2 in the proof of Lemma 12 due to the following:

Lemma 10. *The random variable Z can be expressed as a polynomial of degree 6 in independent uniformly distributed random variables with values -1 and 1 .*

Proof. Consider $C_p = (u, v, w) \in \mathcal{C}$. Let $\epsilon_1^u = -1$ if $\phi(u) = 0$ or 1 and $\epsilon_1^u = 1$, otherwise. Let $\epsilon_2^u = -1$ if $\phi(u) = 0$ or 2 and $\epsilon_2^u = 1$, otherwise. Similarly, we can define $\epsilon_1^v, \epsilon_2^v, \epsilon_1^w, \epsilon_2^w$. Now $\epsilon_1^u \epsilon_2^u$ can be seen as a binary representation of a number from the set $\{0, 1, 2, 3\}$ and $\epsilon_1^u \epsilon_2^u \epsilon_1^v \epsilon_2^v \epsilon_1^w \epsilon_2^w$ can be viewed as a binary representation of a number from the set $\{0, 1, \dots, 63\}$, where -1 plays the role of 0 . Then we can write Z_p as the polynomial

$$\frac{1}{64} \sum_{q=0}^{63} (-1)^{s_q} W_q \cdot (\epsilon_1^u + c_1^{uq})(\epsilon_2^u + c_2^{uq})(\epsilon_1^v + c_1^{vq})(\epsilon_2^v + c_2^{vq})(\epsilon_1^w + c_1^{wq})(\epsilon_2^w + c_2^{wq}),$$

where $c_1^{uq} c_2^{uq} c_1^{vq} c_2^{vq} c_1^{wq} c_2^{wq}$ is the binary representation of q , s_q is the number of digits equal -1 in this representation, and W_q equals the value of Z_p for the case when the binary representations of $\phi(u)$, $\phi(v)$ and $\phi(w)$ are $c_1^{uq} c_2^{uq}$, $c_1^{vq} c_2^{vq}$ and $c_1^{wq} c_2^{wq}$, respectively. The actual values for Z_p for each case are given in the proof of Lemma 9. The above polynomial is of degree 6. It remains to recall that $Z = \sum_{p=1}^m Z_p$. \square

Let us consider the following natural transformation of our key construction introduced in the beginning of this section. Let (V, \mathcal{C}) be an instance of LINEAR ORDERING and (V, \mathcal{B}) , (V, A') and (V, A'') be the associated instances of BETWEENNESS and ACYCLIC SUBDIGRAPH. Let b be the number of pairs of mutually opposite arcs in the directed multigraph $D = (V, A' \cup A'')$ that are deleted by our reduction rule, and let $r = 2(m - b)$. Let t be the number of complete 3-sets of constraints in \mathcal{B} whose deletion from \mathcal{B} eliminates all complete 3-sets of constraints in \mathcal{B} and let $s = m - 3t$.

Lemma 11. *We have $\mathbb{E}[Z^2] \geq \frac{11}{3072}(r + s)$.*

Proof. Let $A = A' \cup A'' = \{a_1, \dots, a_{2m}\}$ and $D = (V, A)$. Fix a function $\phi : V \rightarrow \{0, 1, 2, 3\}$. For a random ϕ -compatible linear ordering π of V , define a binary random variable x_i that takes value one if and only if a_i is satisfied by π . Analogously, define a binary random variable y_i that takes value one if and only if B_i is satisfied by π . Let $X_i = \mathbb{E}[x_i] - 1/2$ for all $i = 1, \dots, 2m$, let $Y_j = \mathbb{E}[y_j] - 1/3$ for all $j = 1, \dots, m$ and let $X = \sum_{i=1}^{2m} X_i$, $Y = \sum_{i=1}^m Y_i$. Recall that b is the number of deleted pairs of mutually opposite arcs from D , and t is the number of complete 3-sets deleted from \mathcal{B} . Assume, without loss of generality, that the remaining arcs are a_1, \dots, a_r and the remaining betweenness constraints are B_1, \dots, B_s . Then $X = \sum_{i=1}^{2m} X_i = \sum_{i=1}^r X_i$, $Y = \sum_{i=1}^m Y_i = \sum_{i=1}^s Y_i$ and, by Lemma 9, $Z = X + Y/2$. Now let ϕ be a random function from V to $\{0, 1, 2, 3\}$. We have the following:

$$\begin{aligned} \mathbb{E}[Z^2] &= \mathbb{E}[X^2 + XY + Y^2/4] = \mathbb{E}[X^2] + \mathbb{E}[Y^2]/4 + \mathbb{E}\left[\left(\sum_{i=1}^r X_i\right)\left(\sum_{j=1}^s Y_j\right)\right] \\ &= \mathbb{E}[X^2] + \mathbb{E}[Y^2]/4 + \sum_{i=1}^r \sum_{j=1}^s \mathbb{E}[X_i Y_j]. \end{aligned}$$

We will show that $\mathbb{E}[X_i Y_j] = 0$ for any pair (i, j) . Let $\phi' : V \rightarrow \{0, 1, 2, 3\}$ be defined as $\phi'(x) = 3 - \phi(x)$ for all x . Let $X_i(\phi)$ be the value of X_i when considering ϕ -compatible orderings and define $X_i(\phi')$, $Y_i(\phi)$ and $Y_i(\phi')$ analogously. From Table 2 we note that $Y_j(\phi) = Y_i(\phi')$, and from Table 3 we note that $X_j(\phi) = -X_i(\phi')$. From $\mathbb{E}[X_i Y_j] = \frac{1}{4^{|V|}} \sum_{\phi} X_i(\phi) Y_j(\phi)$ it follows that

$$2\mathbb{E}[X_i Y_j] = 2 \left[\frac{1}{4^{|V|}} \sum_{\phi} X_i(\phi) Y_j(\phi) \right] = \frac{1}{4^{|V|}} \sum_{\phi} [X_i(\phi) Y_j(\phi) + X_i(\phi') Y_j(\phi')] = 0.$$

Therefore, $\mathbb{E}[Z^2] = \mathbb{E}[X^2] + \mathbb{E}[Y^2]/4$. It follows from Lemmas 4 and 6 that $\mathbb{E}[X^2] \geq r/32$ and $\mathbb{E}[Y^2] \geq \frac{11}{768}s$. We conclude that $\mathbb{E}[Z^2] \geq \frac{11}{3072}(r + s)$. \square

Lemma 12. *There is a constant $c > 0$ such that if $r + s \geq ck^2$, then (V, \mathcal{C}, k) is a “yes”-instance of LINEAR ORDERING-AA.*

Proof. By Lemmas 10 and 2, we have $\mathbb{E}[Z^4] \leq 9^6(\mathbb{E}[Z^2])^2$. As $\mathbb{E}[Z] = 0$, it follows from Lemma 1 that $\mathbb{P}\left(Z > \frac{\sqrt{\mathbb{E}[Z^2]}}{2 \cdot 9^3}\right) > 0$. By Lemma 11, $\mathbb{E}[Z^2] \geq \frac{11}{3072}(r+s)$. Hence, $\mathbb{P}\left(Z > \frac{\sqrt{\frac{11}{3072}(r+s)}}{2 \cdot 9^3}\right) > 0$. Therefore if $r + s \geq ck^2$, where $c = 4 \cdot 9^6 \cdot 3072/11$, then by Lemma 8 (V, \mathcal{C}, k) is a “yes”-instance of LINEAR ORDERING-AA. \square

After we have deleted mutually opposite arcs from D and complete 3-sets of constraints from \mathcal{B} we may assume, by Lemma 12, that D has an arc multiset $A = \{a_1, \dots, a_r\}$ left, with $r = O(k^2)$, and \mathcal{B} now contains $s = O(k^2)$ constraints B_1, \dots, B_s . By Lemma 7, $\text{dev}(V, \mathcal{C}) = \max_{\pi}[(\text{dev}(V, A, \pi) + \text{dev}(V, \mathcal{B}, \pi))/2]$, where the maximum is taken over all linear orderings π of V .

We now create a new instance (V', \mathcal{C}', k) of LINEAR ORDERING-AA as follows. Let ω be a new variable not in V . For every $a_i = (u_i, v_i)$ add the constraints (ω, u_i, v_i) , (u_i, ω, v_i) and (u_i, v_i, ω) to \mathcal{C}' . For every $B_i = (a_i, \{b_i, c_i\})$ add the constraints (b_i, a_i, c_i) and (c_i, a_i, b_i) to \mathcal{C}' . Let V' be the set of variables that appear in some constraint in \mathcal{C}' . Then (V', \mathcal{C}') is an instance of LINEAR ORDERING with $O(k^2)$ variables and constraints. Now the number of constraints in \mathcal{C}' satisfied by any linear ordering α of V' equals the number of arcs in D satisfied by α plus the number of constraints in \mathcal{B} satisfied by α . As the average number of constraints satisfied in (V', \mathcal{C}') equals $(3r + 2s)/6 = r/2 + s/3$, it follows that $\text{dev}(V, \mathcal{C}) = \max_{\pi}[(\text{dev}(V, A, \pi) + \text{dev}(V, \mathcal{B}, \pi))/2] = \text{dev}(V', \mathcal{C}')/2$. Hence, (V', \mathcal{C}', k) is a kernel of LINEAR ORDERING-AA with $O(k^2)$ variables and constraints. We have established the following theorem.

Theorem 3. *LINEAR ORDERING-AA has a kernel with $O(k^2)$ variables and constraints.*

Using Proposition 1 and Theorem 3 we can prove the following:

Theorem 4. *There is a bikernel with $O(k^2)$ variables from Π_i -AA to Π_j -AA for each pair (i, j) such that $0 \leq i \leq 10$ and $0 \leq j \leq 10$ but $j \notin \{2, 7\}$.*

Proof. By Proposition 1, it suffices to prove this theorem for $i = 0$ and $0 \leq j \leq 10$ but $j \notin \{2, 7\}$. The case $j = 0$ follows from Theorem 3. Let us consider the remaining cases.

Part 1: $j = 5$. From the proof of Theorem 3, we know that any instance (V, \mathcal{C}, k) of LINEAR ORDERING-AA can be reduced, in polynomial time, to a mixed instance consisting of an instance $D = (V, A)$ ($|A| = r = O(k^2)$) of ACYCLIC SUBDIGRAPH and an instance (V, \mathcal{B}) ($|\mathcal{B}| = s = O(k^2)$) of BETWEENNESS such that the answer to (V, \mathcal{C}, k) is “yes” if and only if there is a linear ordering of V satisfying, in total, at least $r/2 + s/3 + k$ arcs and constraints of the mixed instance. Let V^* be the set of all variables and vertices in constraints of \mathcal{B} and arcs of A . Observe that $|V^*| = O(k^2)$.

Construct an instance (V', \mathcal{B}', k') of BETWEENNESS-AA as follows. Set $V' = V^* \cup \{y, z\}$ and initialize \mathcal{B}' by setting $\mathcal{B}' = \mathcal{B}$. Add to \mathcal{B}' $(r + s + 1)$ copies of the

constraint $(x, \{y, z\})$ for each $x \in V^*$ and one copy of the constraint $(v, \{u, z\})$ for each arc $(u, v) \in A$. Observe that $|V'| = O(k^2)$. The total number of constraints in the multiset \mathcal{B}' is $p = (|V^*| + 1)(r + s + 1) - 1$ and recall that the average number of constraints satisfied in an instance of BETWEENNESS with p constraints is $p/3$. We may assume that p is divisible by 3 as otherwise we can add one or two more constraints of the type $(x, \{y, z\})$ to \mathcal{B}' . Let $d = (r + s) - \lceil r/2 + s/3 + k \rceil$ and let $k' = \frac{2p}{3} - d$. Observe that the answer to (V', \mathcal{B}', k') is “yes” if and only if there is a linear ordering of V' that falsifies at most d constraints of \mathcal{B}' . Since $d \leq r + s$, to falsify at most d constraints of \mathcal{B}' , a linear ordering α of V' must satisfy all constraints of the form $(x, \{y, z\})$ and at least $r/2 + s/3 + k$ other constraints. Since α must satisfy all constraints of the form $(x, \{y, z\})$, we have $\{\alpha^{-1}(1), \alpha^{-1}(|V'|)\} = \{y, z\}$. Without loss of generality, we may assume that $\alpha^{-1}(|V'|) = z$. Then α satisfies at least $r/2 + s/3 + k$ other constraints if and only if it satisfies at least $r/2 + s/3 + k$ arcs and constraints of the mixed instance. Thus, (V', \mathcal{B}', k') is equivalent to (V, \mathcal{C}, k) , and since k' is bounded by a function of k , we are done.

Part 2: $j = 1$. Denote constraints of Π_1 -AA by $(u < \min\{v, w\})$. Such a constraint is satisfied by a linear ordering α of $\{u, v, w\}$ if and only if $\alpha(u) < \min\{\alpha(v), \alpha(w)\}$. Consider the instance (V', \mathcal{B}', k') built in Part 1. Construct an instance $(V'', \mathcal{C}_1, k_1)$ of Π_1 -AA as follows. Let $V'' = V' \cup \{z'\}$, where $z' \notin V'$. For each constraint $(v, \{u, w\})$ of \mathcal{B}' , let \mathcal{C}_1 have two copies of $(u < \min\{v, w\})$, two copies of $(w < \min\{u, v\})$ and one copy of $(v < \min\{w, z'\})$ and one copy of $(v < \min\{u, z'\})$. Thus, \mathcal{C}_1 has $6p$ constraints and note that the average number of constraints satisfied in an instance of Π_1 -AA with $6p$ constraints is $2p$. Let $k_1 = p - d$, where p and d are defined in Part 1.

Let α be a linear ordering of V'' and assume that α satisfies the maximum number of constraints in \mathcal{C}_1 and this number is at least $2p + k_1 = 3p - d$. We may assume that $\alpha(z') = |V''|$ as moving z' to the last position in the linear ordering will not falsify any constraint of \mathcal{C}_1 . Observe now that if α satisfies $(v, \{u, w\})$, then it satisfies exactly three constraints of \mathcal{C}_1 from the six constraints generated by $(v, \{u, w\})$ and if α falsifies $(v, \{u, w\})$, it satisfies exactly two constraints of \mathcal{C}_1 from the six constraints generated by $(v, \{u, w\})$. Therefore, α satisfies exactly $3t + 2(p - t)$ constraints of \mathcal{C}_1 , where t is the number of constraints in \mathcal{B}' satisfied by α . Hence, $t \geq p - d$.

Now assume that a linear ordering α of V' satisfies at least $p - d$ constraints of \mathcal{B}' . We extend α to V'' by setting $\alpha(z') = |V''|$. Similarly to the above we can show that α satisfies at least $2p + k_1 = 3p - d$ constraints in \mathcal{C}_1 . Thus, (V', \mathcal{C}_1, k_1) is equivalent to (V', \mathcal{B}', k') and, therefore by Part 1, to (V, \mathcal{C}, k) , an instance of LINEAR ORDERING-AA. Clearly, $|V''| = O(k^2)$ and k_1 is bounded by a function of k .

Part 3: $j = 3$. In Part 2, we have proved that for any instance (V, \mathcal{C}, k) of LINEAR ORDERING-AA there is an equivalent instance (V', \mathcal{C}_1, k_1) of Π_1 -AA with $O(k^2)$ variables and distinct constraints (and k_1 is bounded by a function of k). Recall that (V', \mathcal{C}_1, k_1) has $6p$ constraints. Let α be a linear ordering of V' and let α' be the reverse ordering. As in the proof of Case $i = 3$ of Theorem 1, construct from (V', \mathcal{C}_1, k_1) an instance (V', \mathcal{C}_3, k_3) of Π_3 -AA such that \mathcal{C}_3 has $12p$ constraints and at least q constraints of \mathcal{C}_1 are satisfied by α if and only if at least $2q + (|\mathcal{C}_1| - q)$ constraints of \mathcal{C}_3 are satisfied in α' . Let $q = 2p + k_1$ and $k_3 = k_1$. Assume that

(V', \mathcal{C}_1, k_1) is a “yes”-instance certified by α . Then α' satisfies at least $8p + k_3$ constraints of (V', \mathcal{C}_3, k_3) and (V', \mathcal{C}_3, k_3) is a “yes”-instance. Similarly, if (V', \mathcal{C}_3, k_3) is a “yes”-instance, then (V', \mathcal{C}_1, k_1) is a “yes”-instance, too.

Part 4: $j = 4, 8, 9, 10$. For each $j = 4, 8, 9, 10$ the proof is similar to Part 2 and, thus, we will only describe how to transform the instance (V', \mathcal{B}', k') built in Part 1 into an instance (V', \mathcal{C}_i, k') of Π_i -AA for every $i = 4, 8, 9, 10$, and observe how the fact that a constraint B of (V', \mathcal{B}', k') is satisfied or falsified corresponds to the number of satisfied constraints in the instance of Π_i -AA generated by B . Then it is not hard to check that (V', \mathcal{B}', k') and (V', \mathcal{C}_i, k') are equivalent.

Case $j = 4$. Denote constraints of Π_4 -AA by $(u \parallel \{v < w\})$. Such a constraint is Π_4 -satisfied by a linear ordering α of $\{u, v, w\}$ if and only if $\alpha(v) < \alpha(w)$ and $\alpha(u)$ is not between $\alpha(v)$ and $\alpha(w)$. Construct an instance (V', \mathcal{C}_4, k_4) of Π_4 -AA as follows. For each constraint $(v, \{u, w\})$ of \mathcal{B}' , let \mathcal{C}_4 have four constraints: $(u \parallel \{v < w\})$, $(u \parallel \{w < v\})$, $(w \parallel \{u < v\})$ and $(w \parallel \{v < u\})$. It is easy to check that if $(v, \{u, w\})$ is satisfied by a linear ordering α of V' , then two of the four constraints are satisfied by α and if $(v, \{u, w\})$ is falsified by α , then only one of the four constraints is satisfied by α .

Case $j = 8$. Denote constraints of Π_8 -AA by $(v < u < w \text{ or } w < v)$. Such a constraint is satisfied by a linear ordering α of $\{u, v, w\}$ if and only if either $\alpha(v) < \alpha(u) < \alpha(w)$ or $\alpha(w) < \alpha(v)$. For each constraint $(v, \{u, w\})$ of \mathcal{B}' , let \mathcal{C}_8 have two constraints: $(w < v < u \text{ or } u < w)$ and $(u < v < w \text{ or } w < u)$. It is easy to check that if $(v, \{u, w\})$ is satisfied by a linear ordering α of V' , then both constraints generated by $(v, \{u, w\})$ are satisfied by α and if $(v, \{u, w\})$ is falsified by α , then only one of two constraints is satisfied by α .

Case $j = 9$. Denote constraints of Π_9 -AA by $(v \parallel \{u, w\})$. Such a constraint is satisfied by a linear ordering α of $\{u, v, w\}$ if and only if $\alpha(v)$ is not between $\alpha(u)$ and $\alpha(w)$. Construct an instance (V', \mathcal{C}_9, k_9) of Π_9 -AA as follows. For each constraint $(v, \{u, w\})$ of \mathcal{B}' , let \mathcal{C}_9 have two constraints: $(u \parallel \{v, w\})$ and $(w \parallel \{u, v\})$. It is easy to check that if $(v, \{u, w\})$ is satisfied by a linear ordering α of V' , then both constraints generated by $(v, \{u, w\})$ are satisfied by α and if $(v, \{u, w\})$ is falsified by α , then only one of two constraints is satisfied by α .

Case $j = 10$. Denote constraints of Π_{10} -AA by $(\text{not } u < v < w)$. Such a constraint is satisfied by a linear ordering α of $\{u, v, w\}$ if and only if we do not have $\alpha(u) < \alpha(v) < \alpha(w)$. For each constraint $(v, \{u, w\})$ of \mathcal{B}' , let \mathcal{C}_{10} have four constraints: $(\text{not } v < u < w)$, $(\text{not } v < w < u)$, $(\text{not } u < w < v)$ and $(\text{not } w < u < v)$. It is easy to check that if $(v, \{u, w\})$ is satisfied by a linear ordering α of V' , then all four constraints generated by $(v, \{u, w\})$ are satisfied by α and if $(v, \{u, w\})$ is falsified by α , then only three of the four constraints are satisfied by α .

Part 5: $j = 6$. Denote constraints of Π_6 -AA by $(u < v < w \text{ or } w, \{u, v\})$. Such a constraint is satisfied by a linear ordering α of $\{u, v, w\}$ if and only if either $\alpha(u) < \alpha(v) < \alpha(w)$ or $\alpha(w)$ is between $\alpha(u)$ and $\alpha(v)$. Consider the instance

(V', \mathcal{B}', k') built in Part 1. Construct an instance $(V_6, \mathcal{C}_6, k_6)$ of Π_6 -AA as follows.

Let $V_6 = V' \cup \{a, b\}$, where $\{a, b\} \cap V' = \emptyset$. Initiate \mathcal{C}_6 by adding to it, for each $x \in V'$, $6p + 1$ copies of $(x < b < a$ or $a, \{x, b\})$ and $6p + 1$ copies of $(x < a < b$ or $b, \{x, a\})$. For each $(v, \{u, w\}) \in \mathcal{B}'$, add to \mathcal{C}_6 the following constraints: two copies of $(u < w < v$ or $v, \{u, w\})$, two copies of $(w < u < v$ or $v, \{u, w\})$, a copy of $(b < v < u$ or $u, \{v, b\})$, and a copy of $(b < v < w$ or $w, \{b, v\})$. Recall that \mathcal{B}' has p constraints and note that \mathcal{C}_6 has $6p + 2(6p + 1)|V'|$ constraints. Observe that the average number of satisfied constraints, in an instance of Π_6 -AA with $6p + 2(6p + 1)|V'|$ constraints, is $3p + (6p + 1)|V'|$. Let $k_6 = (6p + 1)|V'| + (2p - 3d)$, where d is defined in Part 1.

Then $(V_6, \mathcal{C}_6, k_6)$ is a “yes”-instance if and only if there is a linear ordering α of V_6 that satisfies at least $2(6p + 1)|V'| + (5p - 3d)$ constraints. For α to satisfy so many constraints, it must satisfy all constraints of the forms $(x < b < a$ or $a, \{x, b\})$ and $(x < a < b$ or $b, \{x, a\})$, implying that a and b must be the last two variables in α , and at least $5p - 3d$ constraints generated by \mathcal{B}' . Observe that if α satisfies $(v, \{u, w\}) \in \mathcal{B}'$ then exactly five constraints of \mathcal{C}_6 generated by $(v, \{u, w\})$ are satisfied by α and if α falsifies $(v, \{u, w\}) \in \mathcal{B}'$ then exactly two constraints of \mathcal{C}_6 generated by $(v, \{u, w\})$ are satisfied by α . Thus, α satisfies at least $5p - 3d$ constraints generated by \mathcal{B}' if and only if α satisfies at least $p - d$ constraints of \mathcal{B}' . Therefore, (V', \mathcal{B}', k') and $(V_6, \mathcal{C}_6, k_6)$ are equivalent. \square

Using Theorems 2 and 4 we can prove the following:

Theorem 5. *All ternary Permutation-CSPs parameterized above average have kernels with $O(k^2)$ variables.*

Proof. By Theorem 4, it suffices to prove that the problems Π_j -AA, $j = 2, 7$, have kernels with quadratic number of variables.

Case $j = 2$. Denote constraints of Π_2 -AA by $(u, v < w)$. Such a constraint is satisfied by a linear ordering α of $\{u, v, w\}$ if and only if $\alpha(v) < \alpha(w)$. Consider the instance (V, \mathcal{C}, k) of Π_2 -AA and construct an instance (V, A, k) of ACYCLIC SUBDIGRAPH-AA as follows: if $(u, v < w) \in \mathcal{C}$ then (v, w) is added to A . Clearly, (V, \mathcal{C}, k) and (V, A, k) are equivalent. By Theorem 2, in polynomial time, (V, A, k) can be transformed into an equivalent instance (V', A', k') of ACYCLIC SUBDIGRAPH-AA such that $|V'| = O(k^2)$ and k' is bounded by a function of k (in fact, $k' = k$). As in the proof of Case $i = 2$ of Theorem 1, from (V', A', k') we can construct an equivalent instance (V^*, \mathcal{C}^*, k') of Π_2 -AA such that $|V^*| = |V'| + 1 = O(k^2)$. Observe that (V^*, \mathcal{C}^*, k') is the required kernel.

Case $j = 7$. Denote constraints of Π_7 -AA by $\langle u, v, w \rangle$. Such a constraint is satisfied by a linear ordering α of $\{u, v, w\}$ if and only if either $\alpha(u) < \alpha(v) < \alpha(w)$ or $\alpha(v) < \alpha(w) < \alpha(u)$ or $\alpha(w) < \alpha(u) < \alpha(v)$. Consider the instance (V, \mathcal{C}, k) of Π_7 -AA and construct an instance (V, A, k) of ACYCLIC SUBDIGRAPH-AA as follows: if $\langle u, v, w \rangle \in \mathcal{C}$ then (u, v) , (v, w) and (w, u) are added to A . Let α be a linear ordering of V and observe that if $\langle u, v, w \rangle$ is satisfied by α then exactly two of the three arcs of A generated by $\langle u, v, w \rangle$ are satisfied by α and if $\langle u, v, w \rangle$ is falsified by α then exactly one of the three arcs of A generated by $\langle u, v, w \rangle$ is satisfied by α . Thus, α satisfies at least $|\mathcal{C}|/2 + k$ constraints of \mathcal{C} if and only if α satisfies at least $2(|\mathcal{C}|/2 + k) + (|\mathcal{C}|/2 - k) =$

$3|C|/2 + k = |A|/2 + k$ arcs of A . By Theorem 2, in polynomial time, (V, A, k) can be transformed into an equivalent instance (V', A', k') of ACYCLIC SUBDIGRAPH-AA such that $|V'| = O(k^2)$ and k' is bounded by a function of k (in fact, $k' = k$).

Now construct an instance (V'', C', k') of Π_7 -AA by setting $V'' = V' \cup \{z\}$, where $z \notin V'$, and $C' = \{\langle u, v, z \rangle : (u, v) \in A'\}$. Let α be a linear ordering of V'' satisfying at least $|C'|/2 + k'$ constraints of C' . We may assume that $\alpha(z) = |V''|$ as moving the last element of an ordering to the front of the ordering does not falsify any constraint, and so by repeatedly doing this we will move z' to the last position in our ordering. Thus, α satisfies at least $|A'|/2 + k'$ arcs of A' . Now let α be a linear ordering of V' satisfying at least $|A'|/2 + k'$ arcs of A' . Extend α to V'' by setting $\alpha(z) = |V''|$ and observe that α satisfies at least $|C'|/2 + k'$ constraints in C' . Hence, (V'', C', k') is equivalent to (V', A', k') and, thus, to (V, C, k) implying that (V'', C', k') is a kernel of Π_7 -AA. \square

8 Normal Reduction Rules for Linear Ordering-AA

We call a reduction rule *normal* if it removes a number of constraints which will always have the average number of constraints satisfied no matter what ordering is used. Note that all reduction rules for BETWEENNESS-AA and ACYCLIC SUBDIGRAPH-AA are normal.

Theorem 6 below implies that infinite number of instances of LINEAR ORDERING cannot be reduced by any normal reduction rule, except the one that removes all constraints in the instance. Therefore, no finite set of normal reduction rules can guarantee that one always gets either the empty instance or an instance where one can do better than the average. For both BETWEENNESS-AA and ACYCLIC SUBDIGRAPH-AA we only needed one normal reduction rule to get such a guarantee. This is another indication that LINEAR ORDERING-AA is a more difficult problem.

Let us describe a directed graph G_i with vertex set V_i and a decomposition, C_i , of the arc set of G_i into directed 3-cycles. When $i = 0$ we have $V_0 = \{x_1, x_2, x_3\}$ and $C_0 = \{x_1x_2x_3x_1, x_3x_2x_1x_3\}$. Note that the arc set of G_i is always the set of arcs used in C_i .

When $i > 0$ we will construct G_i , V_i and C_i recursively. So assume that G_{i-1} , V_{i-1} and C_{i-1} have been constructed and let G'_{i-1} be another copy of G_{i-1} on vertex set V'_{i-1} and with decomposition C'_{i-1} . Let $V_i = V_{i-1} \cup V'_{i-1}$ and note that $|V_i| = 2|V_{i-1}|$. Let $c = x_ax_bx_cx_a$ be any directed 3-cycle in C_{i-1} and let $c' = x'_dx'_ex'_fx'_d$ be any directed 3-cycle in C'_{i-1} . Let C_i contain all directed 3-cycles in $C_{i-1} \setminus \{c\}$ and $C'_{i-1} \setminus \{c'\}$ and the following six directed 3-cycles:

$$\begin{array}{lll} c_1 = x_ax_bx'_fx_a, & c_2 = x_bx_cx'_ex_b, & c_3 = x_cx_ax'_dx_c, \\ c_4 = x'_dx'_ex_cx'_d, & c_5 = x'_ex'_fx_bx'_e, & c_6 = x'_fx'_dx_ax'_f. \end{array}$$

A directed graph $D = (V, A)$ is *symmetric* if $(u, v) \in A$ implies $(v, u) \in A$.

Lemma 13. *We have that $|V_i| = 3 \times 2^i$ and that G_i is a symmetric digraph with no parallel arcs for all $i \geq 0$. Furthermore if C_i^* is a proper nonempty subset of C_i then the arcs of C_i^* do not form a symmetric digraph.*

Proof. Since $|V_0| = 3$ and $|V_i| = 2|V_{i-1}|$ we have $|V_i| = 3 \times 2^i$ for all $i \geq 0$. Clearly G_0 is symmetric with no parallel arcs. Assume that G_j is symmetric with no parallel arcs for each $0 \leq j < i$ and consider G_i , $i > 0$. It is not difficult to see that by deleting the arcs in c and c' and adding the arcs in c_1, c_2, \dots, c_6 we obtain a symmetric digraph with no parallel arcs, which completes the proof of the first part of the lemma.

The second part of the lemma clearly holds when $i = 0$, so assume that $i > 0$ and that the second part holds for each $0 \leq j < i$. If $C_i^* \cap \{c_1, c_2, c_3, c_4, c_5, c_6\} = \emptyset$ then we are done by induction as either $C_i^* \cap C_{i-1}$ or $C_i^* \cap C'_{i-1}$ is non-empty and therefore induces a non-symmetric subdigraph.

So we may assume that $C_i^* \cap \{c_1, c_2, c_3, c_4, c_5, c_6\} \neq \emptyset$. Suppose that the arcs of C_i^* form a symmetric digraph. Due to the connection between x_a and x'_f we note that $c_1 \in C_i^*$ if and only if $c_6 \in C_i^*$. Analogously, $c_1 \in C_i^*$ if and only if $c_5 \in C_i^*$ (due to $x'_f x_b$), $c_2 \in C_i^*$ if and only if $c_4 \in C_i^*$ (due to $x'_e x_c$), $c_2 \in C_i^*$ if and only if $c_5 \in C_i^*$ (due to $x'_e x_b$), $c_3 \in C_i^*$ if and only if $c_6 \in C_i^*$ (due to $x'_d x_a$), and $c_3 \in C_i^*$ if and only if $c_4 \in C_i^*$ (due to $x'_d x_c$). Thus, if $C_i^* \cap \{c_1, c_2, c_3, c_4, c_5, c_6\} \neq \emptyset$ and the arcs of C_i^* form a symmetric digraph then we must always have $c_1, c_2, c_3, c_4, c_5, c_6 \in C_i^*$.

As C_i^* is a proper subset of C_i we may without loss of generality assume that there is a directed 3-cycle in $C_{i-1} \setminus \{c\}$ (otherwise it is in $C'_{i-1} \setminus \{c'\}$) which does not belong to C_i^* and by induction the arc set of $(\{c\} \cup C_i^*) \cap C_{i-1}$ does not form a symmetric digraph. Therefore the arcs of C_i^* do not form a symmetric digraph either, a contradiction. This completes the proof of the lemma. \square

For each $i \geq 0$ we construct an instance (V_i, K_i) of LINEAR ORDERING-AA as follows. For every directed 3-cycle in C_i , say $uvwu$, add the following three constraints (u, v, w) , (v, w, u) and (w, u, v) to K_i . Let (V_i, \mathcal{B}_i) be the instance of BETWEENNESS-AA which we associate with (V_i, K_i) in Section 7 and let (V_i, A'_i) and (V_i, A''_i) be the two instances of ACYCLIC SUBDIGRAPH-AA which we also associate with (V_i, K_i) there. By Lemma 7, the following holds for all linear orderings α of V_i :

$$\text{dev}(V_i, K_i, \alpha) = \frac{1}{2} [\text{dev}(V_i, A'_i, \alpha) + \text{dev}(V_i, A''_i, \alpha) + \text{dev}(V_i, \mathcal{B}_i, \alpha)]. \quad (2)$$

Theorem 6. *We have $\text{dev}(V_i, K_i) = 0$ and if K_i^* is a nonempty proper subset of K_i then we can always satisfy more than $|K_i^*|/6$ constraints of K_i^* .*

Proof. As a directed 3-cycle $uvwu$ in C_i gives rise to the betweenness constraints $(v, \{u, w\})$, $(w, \{v, u\})$ and $(u, \{w, v\})$ in \mathcal{B}_i we can only satisfy $|C_i|$ constraints in \mathcal{B}_i . Furthermore, a directed 3-cycle $uvwu$ in C_i gives rise to two copies of the constraints (u, v) , (v, w) and (w, u) in $A'_i \cup A''_i$. Thus, we can think of an arc, uv , in G_i as giving rise to two copies of the acyclic subdigraph constraint (u, v) . As G_i is symmetric this means that every constraint (u, v) can be paired with a constraint (v, u) so we can only satisfy half the constraints in $A'_i \cup A''_i$. As we can only satisfy the average number of constraints in both $A'_i \cup A''_i$ and \mathcal{B}_i , (2) implies that $\text{dev}(V_i, K_i) = 0$, which proves the first part of the lemma.

For the sake of contradiction assume that K_i^* is a nonempty proper subset of K_i and that $\text{dev}(V_i, K_i^*) = 0$. Let (V_i, \mathcal{B}_i^*) be the instance of BETWEENNESS-AA which we associate with (V_i, K_i^*) in Section 7 and let (V_i, A_i^*) and (V_i, A_i^{**}) be the two instances of ACYCLIC SUBDIGRAPH-AA which are also associated with (V_i, K_i^*) . Let Z , Y and X be the random variables associated with (V_i, K_i^*) , (V_i, \mathcal{B}_i^*) and

$(V_i, A_i^* \cup A_i^{**})$, respectively. Note that $\text{dev}(V_i, K_i^*) = 0$ is equivalent to $\mathbb{E}[Z^2] = 0$, which by the proof of Lemma 11 implies that $\mathbb{E}[X^2] = 0$ and $\mathbb{E}[Y^2] = 0$. Observe that by Lemma 4 this implies that if $(u, \{v, w\}) \in \mathcal{B}_i^*$ then $(w, \{v, u\}), (v, \{u, w\}) \in \mathcal{B}_i^*$. So, if $(u, v, w) \in K_i^*$, then $(v, w, u), (w, u, v) \in K_i^*$. Therefore, K_i^* can be thought of as being obtained from a proper subset, C_i^* , of the directed 3-cycles C_i . Observe that by Lemma 13 some arc (u, v) belongs to a directed 3-cycle in C_i^* , but the arc vu does not belong to such a directed 3-cycle. However, this implies that $(u, v) \in A_i^* \cup A_i^{**}$, but $(v, u) \notin A_i^* \cup A_i^{**}$. Thus, $\mathbb{E}[X^2] > 0$ by Lemma 6. This contradiction completes the proof. \square

9 Further Research

It is natural and easy to extend the definition of Π -AA to a fixed arity $r > 3$. Similar to Proposition 1, it is easy to prove that, for each fixed r every Π -AA can be reduced to Π_0 -AA, where Π_0 is LINEAR ORDERING of arity r . However, it appears technically very difficult to extend results obtained for arities $r = 2$ and 3 to $r > 3$. We conjecture that for each fixed r all problems Π -AA are fixed-parameter tractable.

We have parameterized LINEAR ORDERING of any arity r using the average as a tight lower bound. Similarly, we can parameterize LINEAR ORDERING below a tight upper bound and the number of constraints m can be set as a tight upper bound. So, the problem is whether there is a bijection $\alpha : V \rightarrow [n]$ which satisfies at least $m - k$ constraints of an instance (V, \mathcal{C}) of LINEAR ORDERING, where k is the parameter. It is easy to show that for $k = 0$ the problem is polynomial-time solvable, but it seems to be a difficult question to determine the parameterized complexity of this problem for any arity $r \geq 2$.

Note that for arity $r = 2$ the corresponding problem is DIRECTED FEEDBACK ARC SET parameterized below the number m of arcs in a given directed graph. The parameterized complexity of the last problem was an open question for many years [24] and, only in 2008, Chen et al. [7] proved that the problem is fixed-parameter tractable. (It is still unknown whether the last problems admits a polynomial-size kernel.) For every fixed arity $r \geq 3$, the parameterized complexity of LINEAR ORDERING parameterized below m is unknown.

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